

A Note On Spectral Clustering



Pavel Kolev and Kurt Mehlhorn



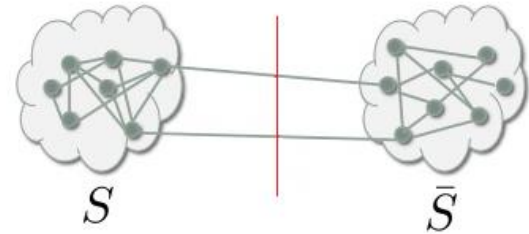
European Symposia on Algorithms'16

Outline

- **Problem Formulation**
 - Algorithmic Tools
- **Our Contribution**
 - Structural Result
 - Algorithmic Result
 - Proof Overview
- **Summary**

k-way Partitioning

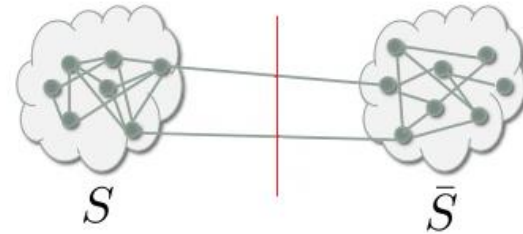
- **Def.** A **cluster** is a subset $S \subseteq V$ with small **conductance**



$$\phi(S) = \frac{|E(S, \bar{S})|}{\mu(S)}, \text{ where the volume } \mu(S) = \sum_{v \in S} \deg(v).$$

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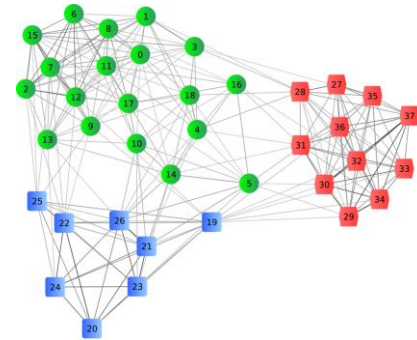
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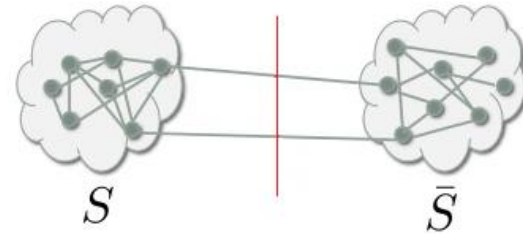
- **Def.** The order k conductance constant

$$\rho(k) = \min_{\text{partition } (P_1, \dots, P_k)} \max_{i \in [1:k]} \phi(P_i)$$



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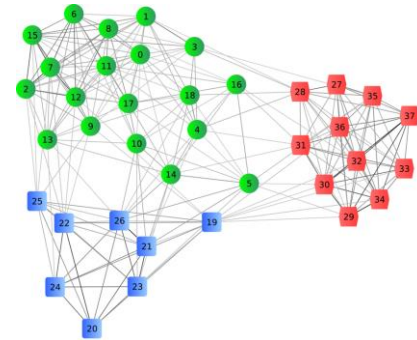
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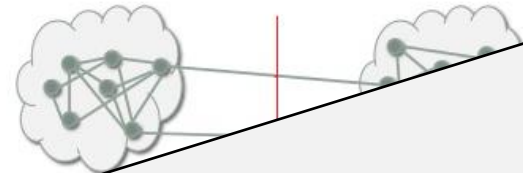


- **Goal:** Find an **approximate** k -way partition w.r.t $\rho(k)$.

k-way Partitioning

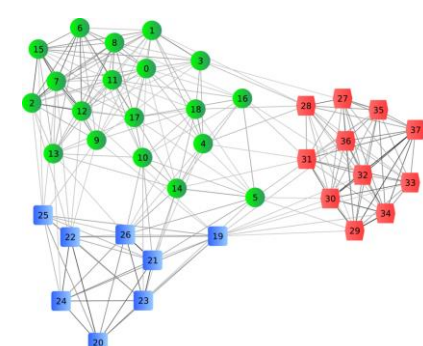
- **Def.** A **cluster** is a subset $S \subseteq V$ with small **conductance**

$$\phi(S) = \frac{|E(S, \bar{S})|}{\mu(S)}, \text{ where the } \mu(S) \text{ is the volume of } S.$$



Here, we analyze rigorously the **Standard Spectral Clustering Paradigm**.

$$\min_{\text{partition } (P_1, \dots, P_k)} \max_{i \in [1:k]} \phi(P_i)$$



- **Goal:** Find an **approximate** k -way partition w.r.t $\rho(k)$.

Standard Spectral Clustering Paradigm

Input: $G = (V, E)$, $3 \leq k \ll n$ and $\epsilon \in (0,1)$.

Output: An *approximate* k -way partition of V .

Andrew Ng et al [NIPS'02]:

1. Computes an *approximate* Spectral Embedding ($F: V \mapsto R^k$) using the Power Method.
- 2) Run a k -means clustering algorithm to compute an *approximate* k -way partition of $\{F(v)\}_{v \in V}$.

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Spectral Graph Theory

- The **normalized Laplacian** matrix \mathcal{L} has **eigenvalues**

$$0 = \lambda_1 \leq \dots \leq \lambda_k \leq \lambda_{k+1} \leq \dots \leq \lambda_n \leq 2.$$

- **Fact.** A graph has exactly k **connected component** iff

$$0 = \lambda_k < \lambda_{k+1}.$$

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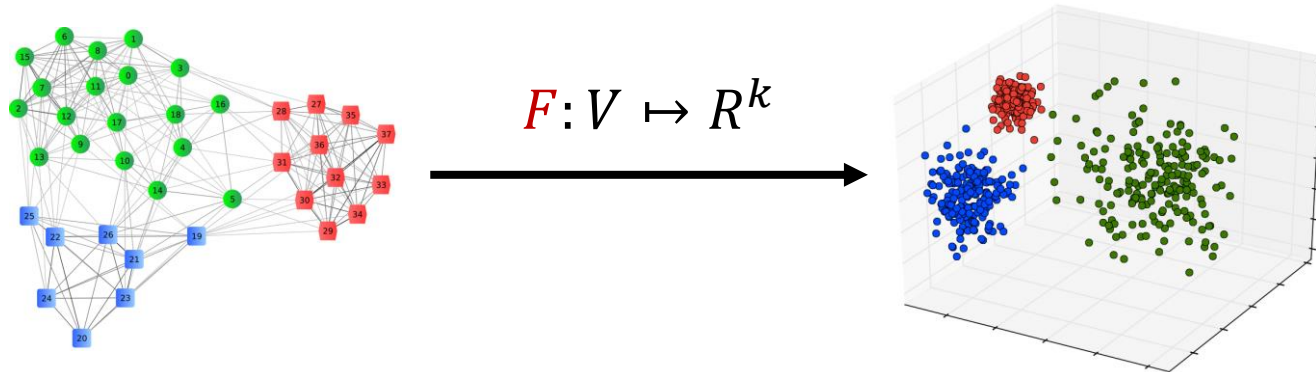
- **Trevisan et al. [STOC'12, SODA'14]** proved a **robust** version

$$\lambda_k/2 \leq \rho(k) \leq O(k^3)\sqrt{\lambda_k}.$$

$(\rho(k))$ is **NP-hard** and λ_k is in **P**) \rightarrow **approx. scheme!**

Exact Spectral Embedding

- $U_k = (v_1, v_2, \dots, v_k) \in R^{V \times k}$ - the bottom k eigenvectors of \mathcal{L}

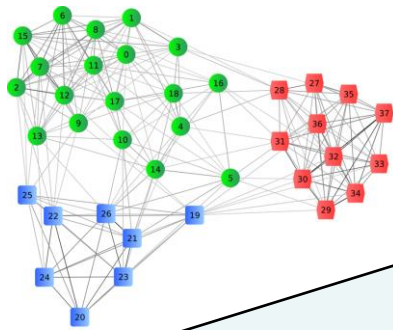


- Normalized Spectral Embedding:

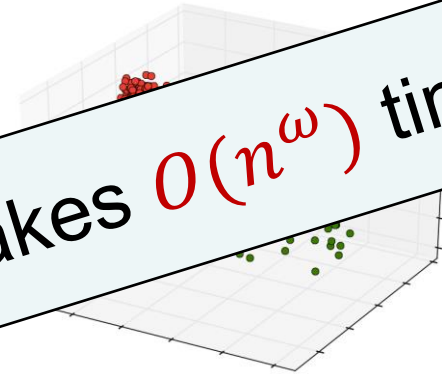
$$F(v) = \frac{1}{\sqrt{\deg(v)}} U_k(v, :), \text{ for every } v \in V.$$

Exact Spectral Embedding

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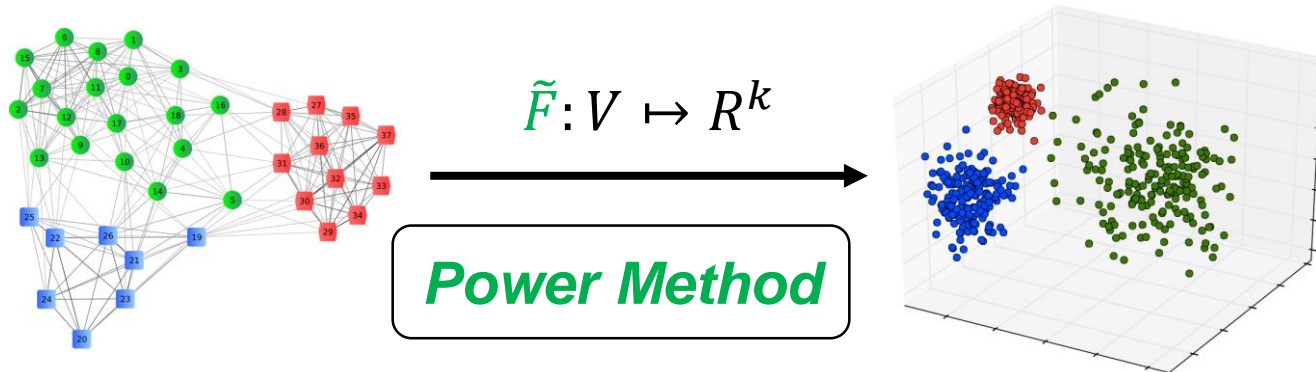
The **exact** computation of U_k takes $O(n^\omega)$ time!

- **Fast Spectral Embedding:**

$$F(v) = \frac{1}{\sqrt{\deg(v)}} U_k(v, :), \text{ for every } v \in V.$$

Approximate Spectral Embedding

- $\tilde{U}_k \in R^{V \times k}$ approximation of the bottom k eigenvectors of \mathcal{L}

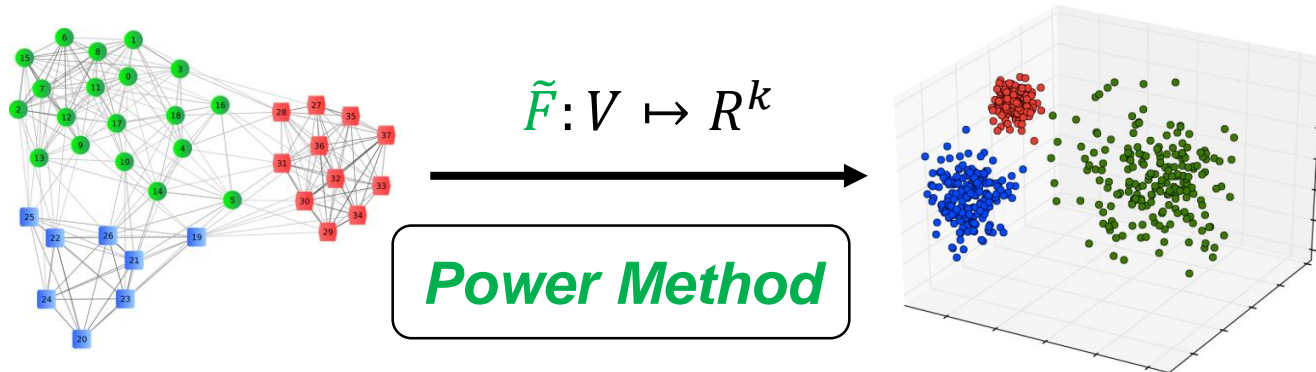


- Approximate Normalized Spectral Embedding:

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Approximate Spectral Embedding

- $\tilde{U}_k \in R^{V \times k}$ approximation of the bottom k eigenvectors of \mathcal{L}



- **Approximate Normalized Spectral Embedding:**

Point Sets:

$$\tilde{\mathcal{X}}_E = \{\text{deg}(v) \text{ many copies of } \tilde{F}(v) \mid v \in V\}.$$

$$\tilde{\mathcal{X}}_V = \{\tilde{F}(v) \mid v \in V\}.$$

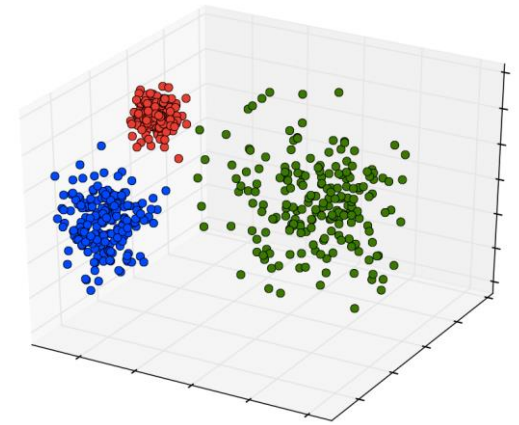
k -means Clustering

Input: $\mathcal{X} = \{p_1, \dots, p_n\}$ with $p_i \in R^k$.

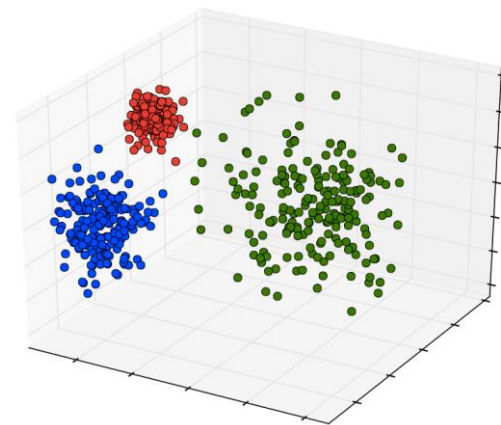
Output: k -way partition of \mathcal{X} such that

$$(A_1^*, \dots, A_k^*) = \underset{\text{partition } (X_1, \dots, X_k) \text{ of } \mathcal{X}}{\operatorname{argmin}} \sum_{i=1}^k \sum_{p \in X_i} \|p - c_i\|^2 ,$$

where c_i is the center of X_i .



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Def. The optimal k -means cost is

$$\Delta_k(\mathcal{X}) = \operatorname{cost}(A_1^*, \dots, A_k^*).$$

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Structural Result

- **Peng et al. [COLT'15]**

$$\Upsilon := \lambda_{k+1} / \rho(k) \geq \Omega(k^3)$$

$$\rho(k) = \max_{i \in [1:k]} \phi(P_i)$$

- **Our Result**

$$\Psi := \lambda_{k+1} / \rho_{\text{avr}}(k) \geq \Omega(k^3)$$

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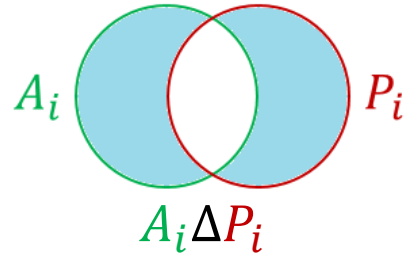
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- $\text{cost}(A_1, \dots, A_k) \leq \gamma \cdot \Delta_k(\tilde{\mathcal{X}}_E)$ for $\gamma \geq 1$.

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How to find such k -way partition (A_1, \dots, A_k) ?

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Algorithmic Result

- Peng et al. [COLT'15]

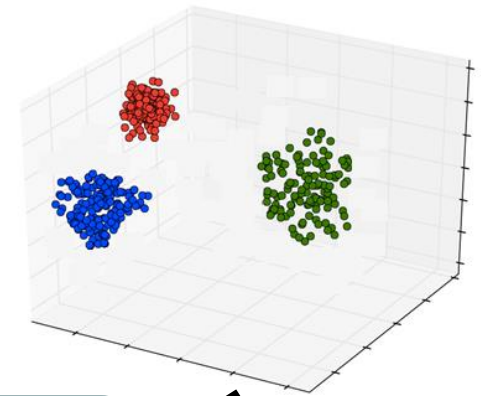
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more restrictive by

$\Omega(k^2)$ -factor

Concentration

Heat Kernel and
Local Sensitive Hashing



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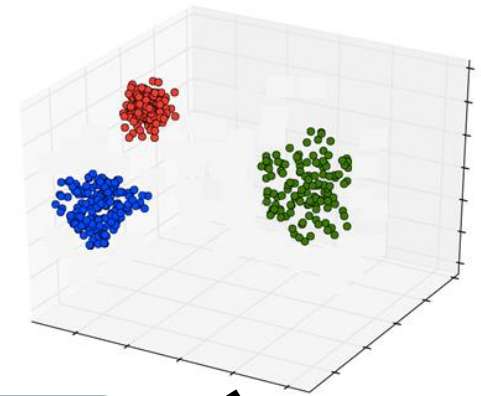
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- Our Result

$$\Psi := \lambda_{k+1} / \rho_{\text{avr}}(k) \geq \Omega(k^3)$$

$$\text{and } \Delta_k(\mathcal{X}_V) \geq n^{-O(1)}$$

Approx. Spectral Embedding
and k-means Clustering

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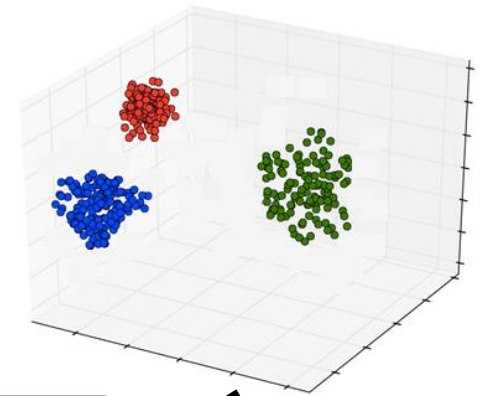
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This is the 1st rigorous *algorithmic* analysis of the
Standard Spectral Clustering Paradigm!

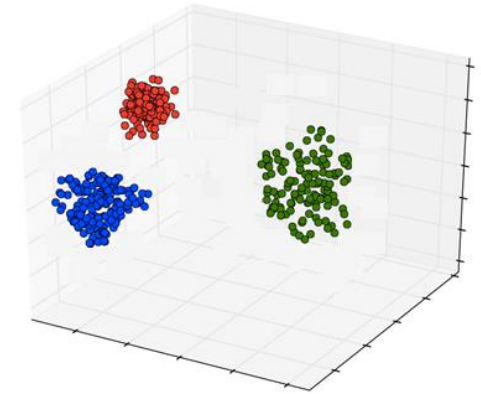
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constant = 10^5

Concentration



Heat Kernel and Local Sensitive Hashing

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constant = $10^7 / \epsilon_0$

Approx. Spectral Embedding
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$\epsilon_0 = 6/10^7$ is Ostrovsky et al's [FOCS'13]
k-means alg. constant (is not optimized!)

Algorithmic Result

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Heat Kernel and
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Heat Kernel and
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Runtime: $O(m \log^c n)$

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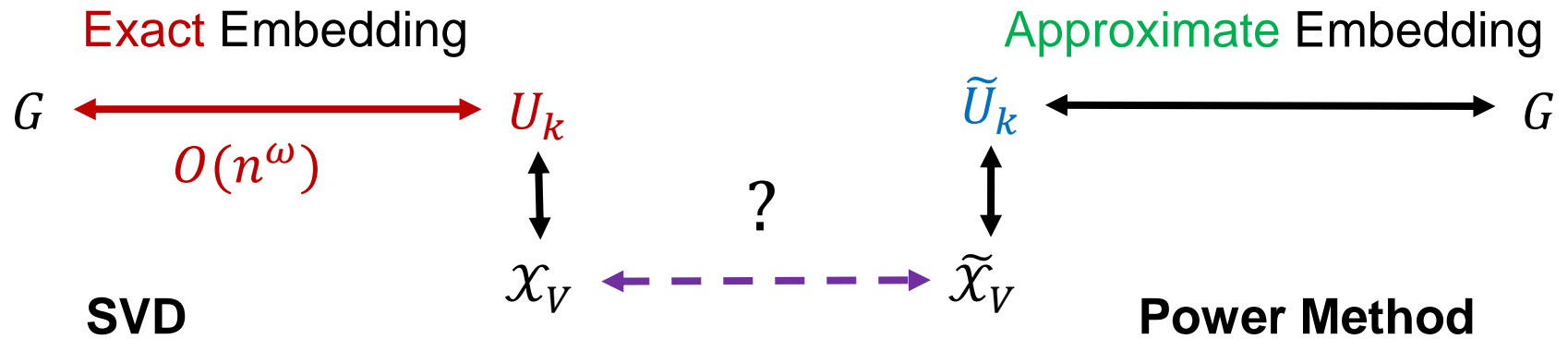
Approx. Spectral Embedding
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Runtime: $O\left(m \left(k^2 + \frac{\ln n}{\lambda_{k+1}}\right)\right)$

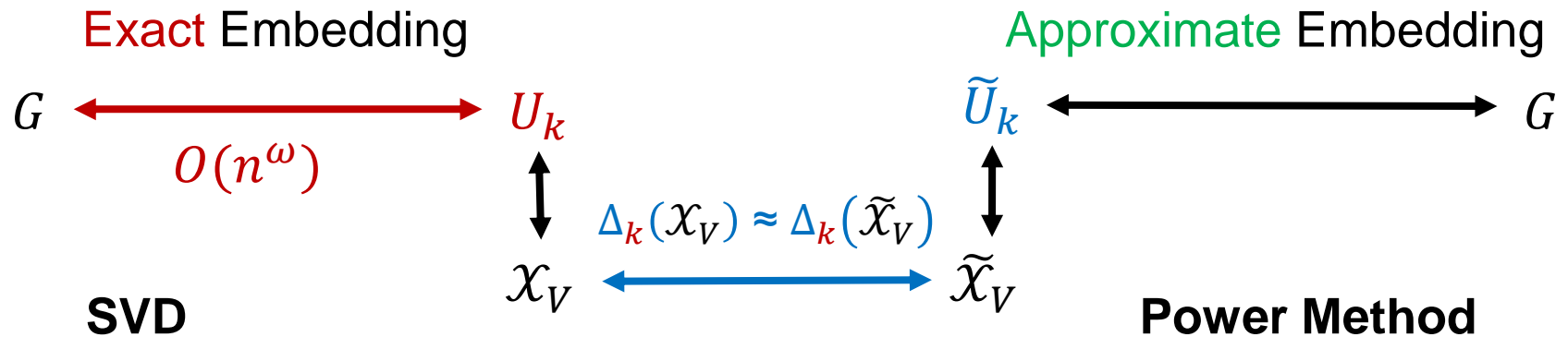
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Proof Overview



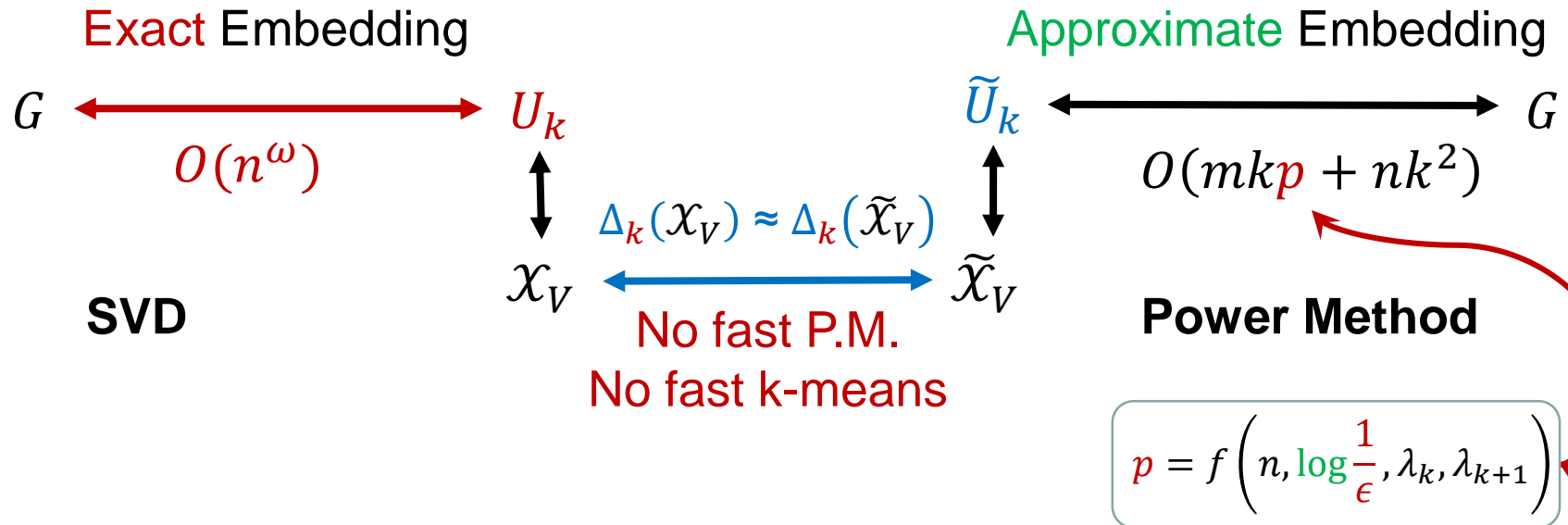
Boutsidos et al [ICML'15] Let (A_1, \dots, A_k) be a partition such that

$$\text{cost}(A_1, \dots, A_k) \leq (1 + \gamma)\Delta_k(\tilde{\mathcal{X}}_V)$$

then

$$\text{cost}(A_1, \dots, A_k) \leq (1 + 4\epsilon)(1 + \gamma)\Delta_k(\mathcal{X}_V) + 4\epsilon^2.$$

Proof Overview



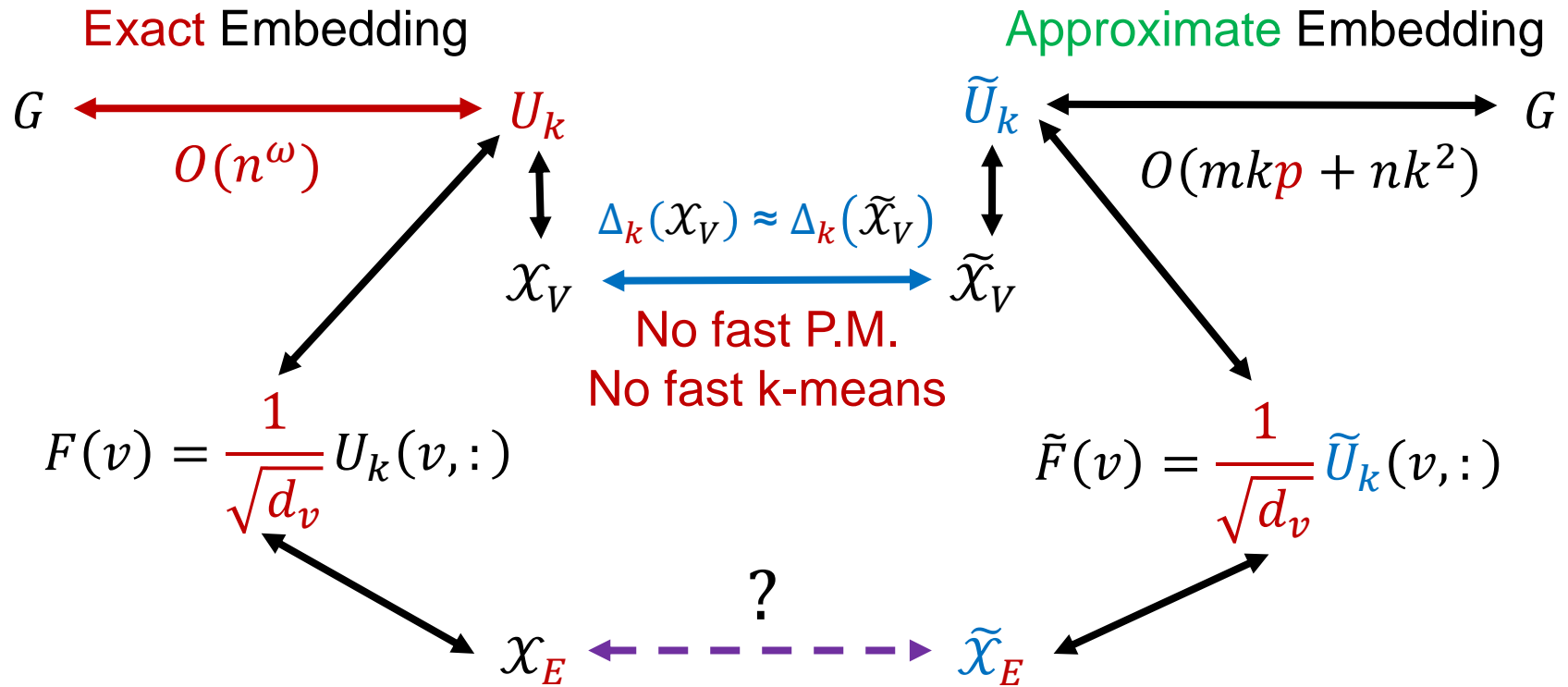
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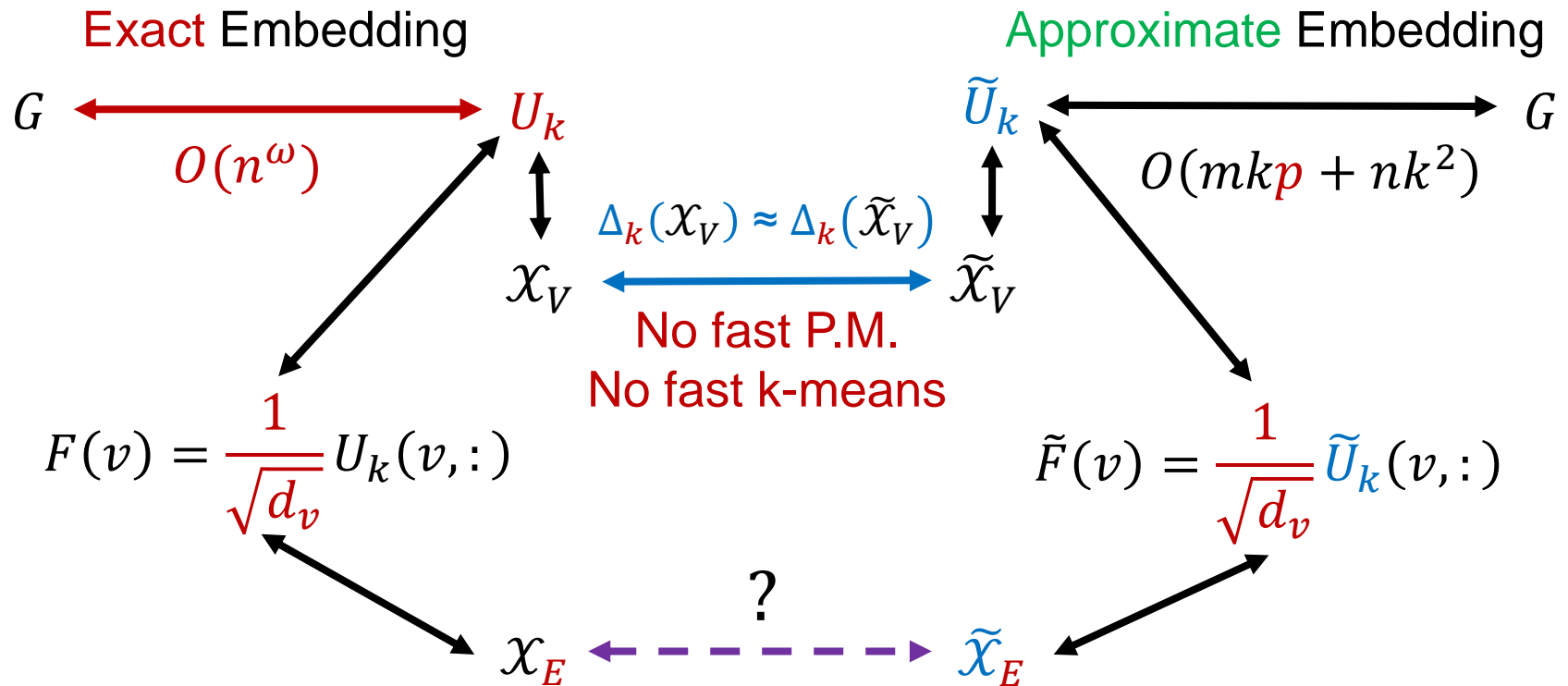
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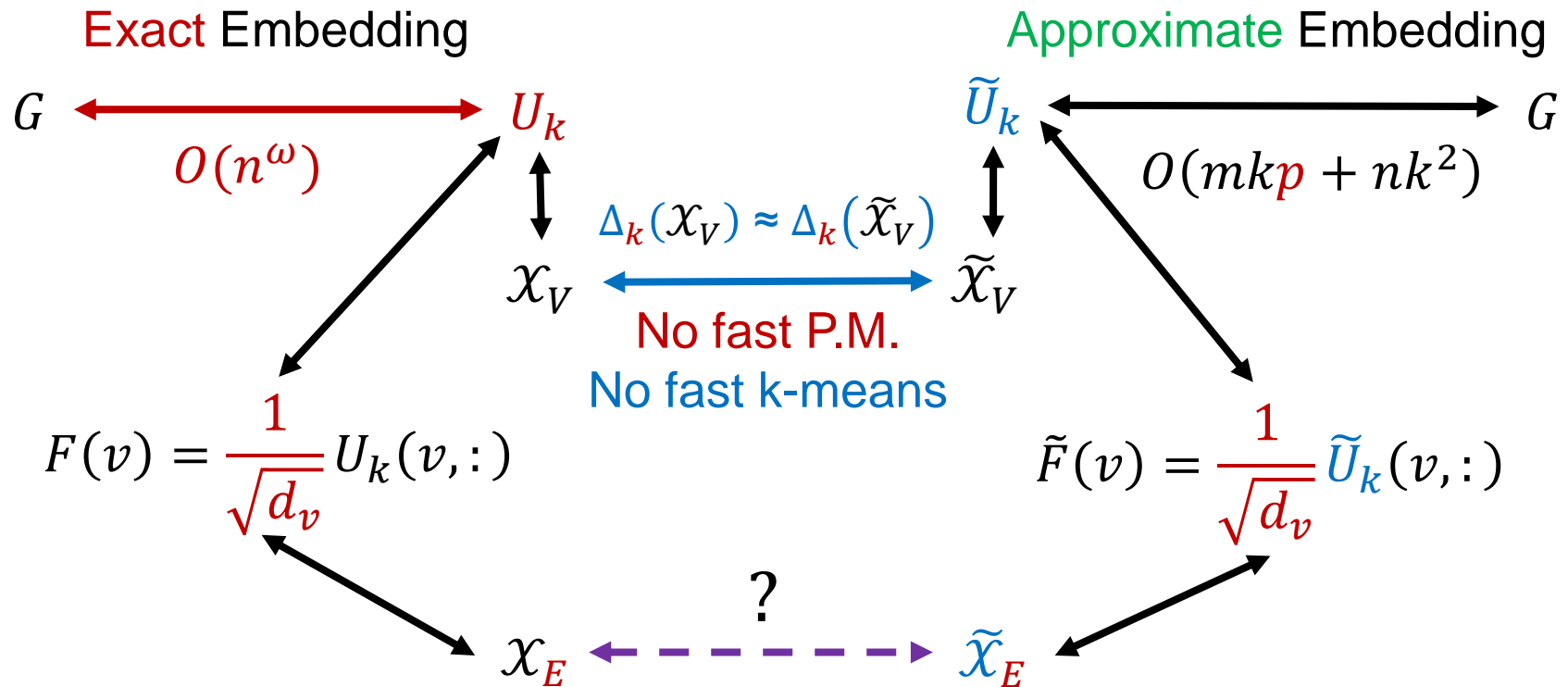
Proof Overview



Questions:

1. Find an efficient k -means clustering algorithm for \tilde{x}_E ?
2. Extend Boutsidos et al's [ICML'15] analysis?

Proof Overview

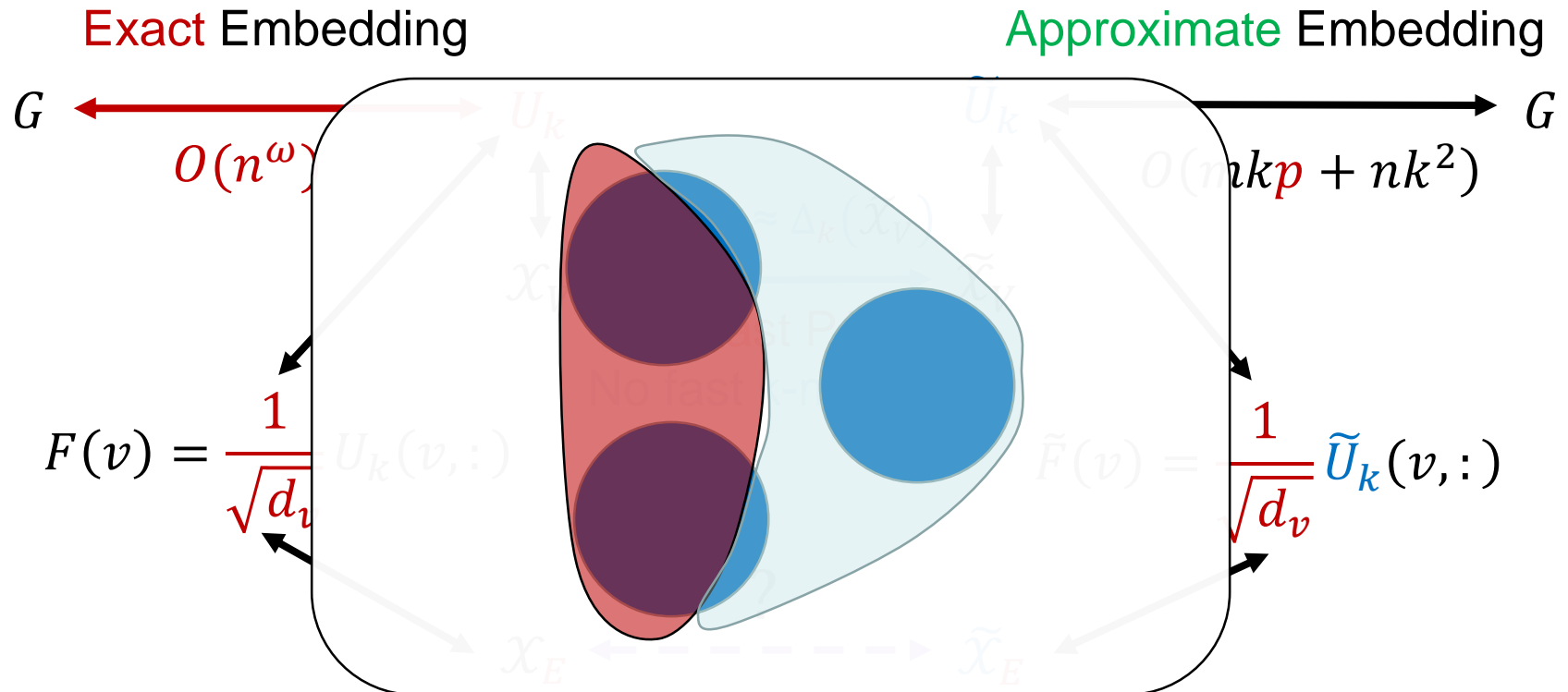


Ostrovsky et al's [FOCS'13] gave an *approximate* k-means algorithm

with fast runtime $O(mk^2)$, but requires $\Delta_k(\mathcal{X}) \leq \epsilon_0^2 \cdot \Delta_{k-1}(\mathcal{X})$

where $\epsilon_0 = 6/10^7$.

Proof Overview



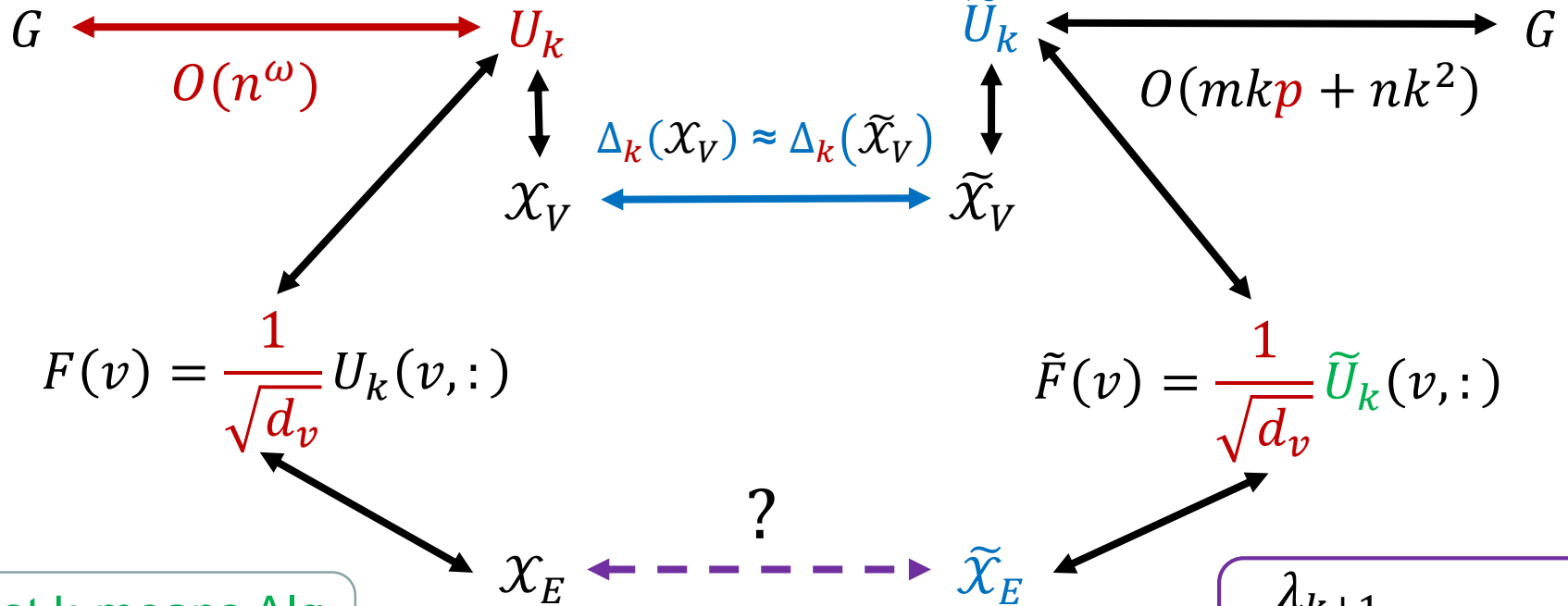
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Proof Overview

Exact Embedding

Approximate Embedding



Fast k-means Alg.
runtime: $O(mk^2)$

$$\frac{\lambda_{k+1}}{\rho_{\text{avr}}(k)} = \Omega(k^3)$$

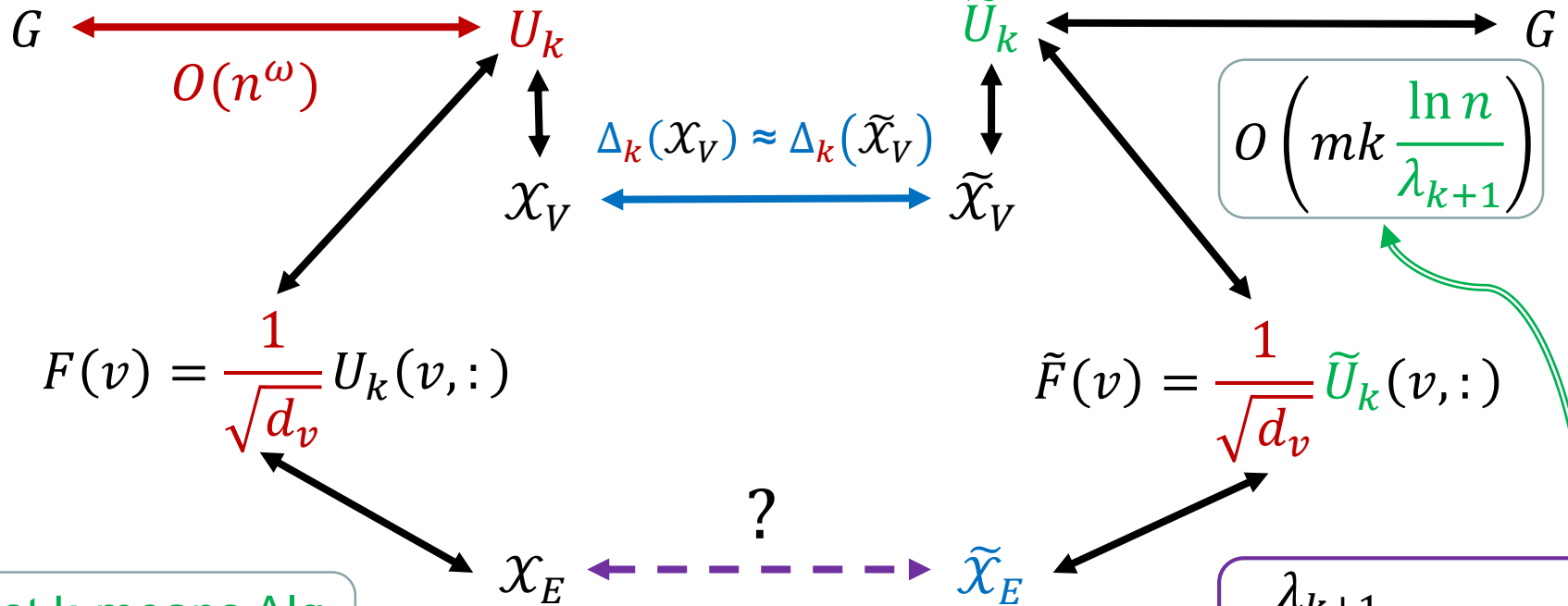
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Proof Overview

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$$O\left(mk \frac{\ln n}{\lambda_{k+1}}\right)$$

Fast k-means Alg.
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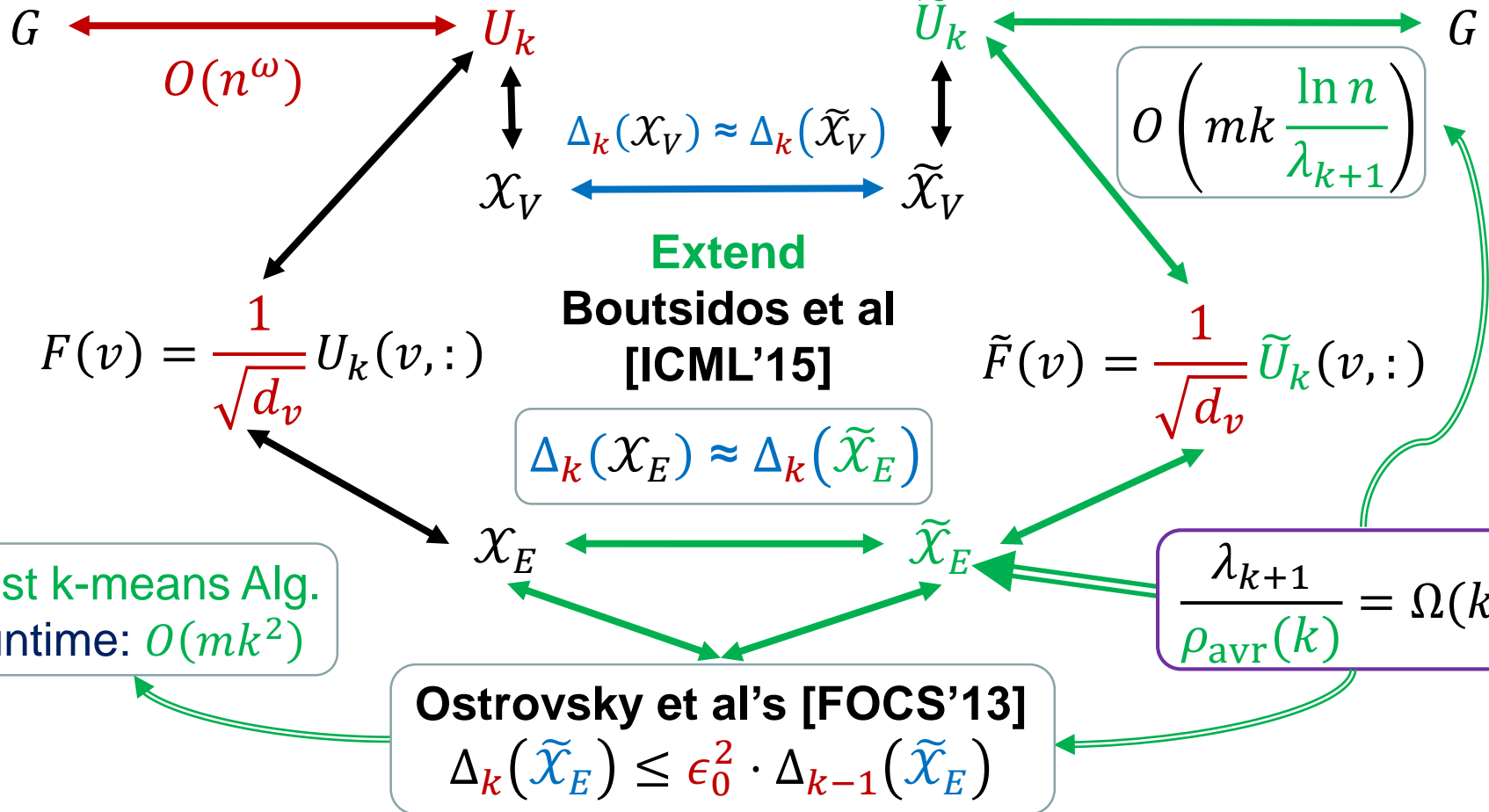
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Proof Sketch (Overview)

Exact Embedding

Approximate Embedding



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Summary

- *We proved rigorously that*

the Standard Spectral Clustering Paradigm

efficiently computes a k -way partition

under asymptotically less restrictive gap assumption.

Open Problems

- Show that the **SSCP** has a **good** behavior on **small** graphs.

Our approach **fails** due to **large** constants in $\Psi \geq \Omega(k^3)$:

- $10^7/\epsilon_0$ - **Ostrovsky et al.** (is not optimized)

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Our approach **fails** due to **large** constants in $\Psi \geq \Omega(k^3)$:

- $10^7/\epsilon_0$ - **Ostrovsky et al.** (is not optimized)

$$\Delta_k(\mathcal{X}) \leq \epsilon_0^2 \cdot \Delta_{k-1}(\mathcal{X}), \text{ where } \epsilon_0 = 6/10^7.$$

- Can we obtain a **multiplicative conductance** guarantee:

$$\phi(A_i) \leq (1 + \gamma/\Psi k) \cdot \phi(P_i) + \gamma/\Psi k.$$

Thank you!